

Math 373: Mathematical Programming and Optimization I

Fall, 2024 Assignment 2

October 14, due October 26

1. Let \mathbf{x} be an element of the polyhedron $P = \{\mathbf{x} \in \mathbb{R}^n : \mathbf{A}\mathbf{x} = \mathbf{b}, \mathbf{x} \geq \mathbf{0}\}$. Prove that a vector $\mathbf{d} \in \mathbb{R}^n$ is a feasible direction at \mathbf{x} if and only if $\mathbf{A}\mathbf{d} = \mathbf{0}$ and $d_i \geq 0$ for every i such that $x_i = 0$. 5

If \mathbf{d} is a feasible direction, then $\mathbf{x} + t\mathbf{d} \in P$ for some positive scalar t . That is, $\mathbf{A}(\mathbf{x} + t\mathbf{d}) = \mathbf{b}$ and $\mathbf{x} + t\mathbf{d} \geq \mathbf{0}$. Then $t\mathbf{A}\mathbf{d} = \mathbf{A}(\mathbf{x} + t\mathbf{d}) - \mathbf{A}\mathbf{x} = \mathbf{b} - \mathbf{b} = \mathbf{0}$, so that $\mathbf{A}\mathbf{d} = \mathbf{0}$. Moreover, for each zero component x_i , the condition $x_i + td_i \geq 0$ reduces to $d_i \geq 0$.

Conversely, if there exists a direction such that $\mathbf{A}\mathbf{d} = \mathbf{0}$ and $d_i \geq 0$ for every i such that $x_i = 0$, then $\mathbf{A}(\mathbf{x} + t\mathbf{d}) = \mathbf{A}\mathbf{x} + t\mathbf{A}\mathbf{d} = \mathbf{b} + \mathbf{0} = \mathbf{b}$ for every real t . We are also given that $x_i + td_i = td_i \geq 0$ for every component i such that $x_i = 0$. For those i for which $x_i > 0$, then unless $d_i = 0$ (in which case $x_i + td_i > 0$ for every t) let us enforce $t \leq x_i/|d_i| > 0$, so that $x_i + td_i \geq |td_i| + td_i \geq 0$. On choosing $t^* = \min_{x_i > 0, d_i \neq 0} x_i/|d_i| > 0$,

we thus see that $\mathbf{x} + t^*\mathbf{d} \geq \mathbf{0}$. Hence \mathbf{d} is a feasible direction at \mathbf{x} .

2. Let \mathbf{x} be a basic feasible solution of a linear programming problem Π written in standard form, with associated basis matrix \mathbf{B} and set of nonbasic indices N . Let \mathbf{y} be any feasible solution to Π and consider the difference vector $\mathbf{d} = \mathbf{y} - \mathbf{x}$.

(a) Prove that $d_j \geq 0$ for every $j \in N$. 1

For any feasible solution \mathbf{y} we have $\mathbf{y} \geq \mathbf{0}$. Since \mathbf{x} is a basic feasible solution, we know for each $j \in N$ that $x_j = 0$ and hence $d_j = y_j - x_j \geq 0$.

(b) If $d_j = 0$ for every $j \in N$, prove that $\mathbf{y} = \mathbf{x}$. 1

This would imply that

$$\mathbf{0} = \mathbf{A}\mathbf{y} - \mathbf{A}\mathbf{x} = \mathbf{A}\mathbf{d} = \mathbf{B}\mathbf{d}_B + \sum_{j \in N} \mathbf{A}_j d_j = \mathbf{B}\mathbf{d}_B.$$

The linear independence of the columns of \mathbf{B} then implies that $\mathbf{d}_B = \mathbf{0}$ and hence $\mathbf{d} = \mathbf{0}$, so that $\mathbf{y} = \mathbf{x}$.

(c) If the reduced cost \bar{c}_j of every nonbasic variable x_j is positive, use parts (a) and (b) to prove that \mathbf{x} is the unique optimal solution to Π . 2

Recall that \bar{c}_j is the rate of change along the j th simplex direction. That is, the change in cost on moving from \mathbf{x} to \mathbf{y} is

$$\mathbf{c}^\top \mathbf{y} - \mathbf{c}^\top \mathbf{x} = \mathbf{c}^\top \mathbf{d} = \mathbf{c}_B^\top \mathbf{d}_B + \sum_{j \in N} c_j d_j = \sum_{j \in N} (c_j - \mathbf{c}_B^\top \mathbf{B}^{-1} \mathbf{A}_j) d_j = \sum_{j \in N} \bar{c}_j d_j.$$

We know from part (a) that $d_j \geq 0$. Moreover, if $\mathbf{y} \neq \mathbf{x}$, we know from part (b) that $d_j > 0$ for some $j \in N$. Given $\bar{c}_j > 0$ for each $j \in N$, we see that

$$\mathbf{c}^T \mathbf{y} - \mathbf{c}^T \mathbf{x} = \sum_{j \in N} \bar{c}_j d_j > 0.$$

Since this holds for every feasible vector $\mathbf{y} \neq \mathbf{x}$, we see that \mathbf{x} is the unique optimal solution.

(d) Suppose that \mathbf{x} is a nondegenerate optimal solution to Π . If the reduced cost \bar{c}_j of some nonbasic variable x_j is zero at \mathbf{x} , prove that Π does not have a unique optimal solution. 1

Let \mathbf{d}' be the j th simplex direction. Since \mathbf{x} is nondegenerate, we know that the solution $\mathbf{y} = \mathbf{x} + t\mathbf{d}'$ is feasible for some $t > 0$. From the definition of the j th simplex direction, we see that

$$\mathbf{c}^T \mathbf{y} - \mathbf{c}^T \mathbf{x} = t\bar{c}_j d'_j = 0.$$

That is, \mathbf{y} is a distinct feasible solution with the same optimal cost as \mathbf{x} .