

Chromatic Polynomials

Consider the following question: Given a graph G and k colors, how many different ways are there to color G with these colors?

Definition and examples

Definition 1. (Chromatic Polynomial) Let G be a simple graph of order n with vertices marked v_1, \dots, v_n . For every $k \in \mathbb{N}$, let $P_G(k)$ be the number of ways to color the graph G with k colors so that any two vertices connected by an edge are colored differently. $P_G(k)$ is called the chromatic polynomial of the graph G .

Remark 2. In other words, $P_G(k)$ is the number of functions σ from $V = \{v_1, \dots, v_n\}$ to $\{1, 2, \dots, k\}$ so that $\sigma(v_i) \neq \sigma(v_j)$ if $\{v_i, v_j\} \in E$.

Remark 3. It is clear that $P_G(k) = 0$ when $k < \chi(G)$. Thus any root to $P_G(k)$ gives a lower bound for $\chi(G)$. The original motivation of Birkhoff to introduce the idea of chromatic polynomial in 1912 is to solve the four-color problem as follows:

- i. Characterize chromatic polynomials for planar graphs;
- ii. Characterize the roots of these polynomials.

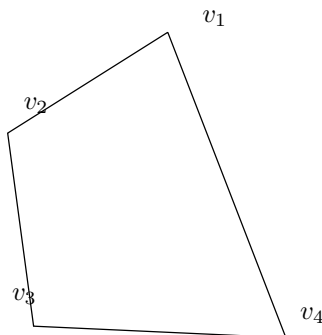
Unfortunately, to this day we are still stuck at i.

Example 4. Let N_n be the null graph of order n . We have $P_{N_n}(k) = k^n$.

Example 5. Let K_n be the complete graph of order n . We have $P_{K_n}(k) = k(k-1) \cdots (k-(n-1))$, as there are k colors available for v_1 , then $k-1$ colors available for v_2, \dots , $k-(n-1)$ colors available for v_n .

Exercise 1. Let G be any simple graph of order n . Then $P_{K_n}(k) \leq P_G(k) \leq P_{N_n}(k)$.

Example 6. ¹We try to calculate $P_G(k)$ for the graph



1. Note that this is a motivational example. Later we will discuss better ways of calculating $P_G(k)$ for graphs.

First we notice that $P_G(1)=0$ and $P_G(2)=2$. In general, for $k > 2$, let's assume that one of the colors is red. Then $P_G(k) = k R(k)$ where $R(k)$ is the number of ways coloring the graph with v_1 colored red. To calculate $R(k)$, we divide into two cases (and then apply the sum rule).

- v_3 is also colored red. In this case v_2 has $k - 1$ choices and v_4 too. Thus we have $(k - 1)^2$.
- v_3 is not colored red. In this case v_3 has $k - 1$ choices and v_2, v_4 have $(k - 2)$ choices each. Consequently we have $(k - 1)(k - 2)^2$.

Thus overall we have

$$P_G(k) = k [(k - 1)^2 + (k - 1)(k - 2)^2] = k(k - 1)(k^2 - 3k + 3) = k^4 - 4k^3 + 6k^2 - 3k. \quad (1)$$

Remark 7. Note that we can become more confident of our calculation of $P_G(k)$ by checking $P_G(k)$ for small values of k :

$$P_G(0) = 0; \quad P_G(1) = 1; \quad P_G(2) = 2. \quad (2)$$

Properties

Theorem 8. (Deletion-Contraction) Let $G = (V, E)$ be a simple graph. Let $e = \{a, b\} \in E$. Let G_D be the graph obtained from G by deleting the edge e , and let G_C be the graph obtained from G by identifying the two vertices a, b . Then

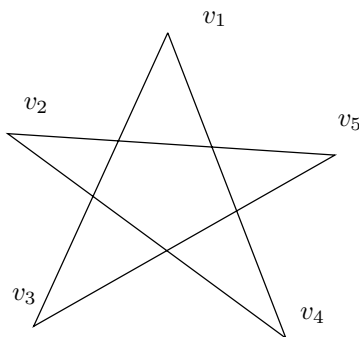
$$P_G(k) = P_{G_D}(k) - P_{G_C}(k). \quad (3)$$

Proof. It is clear that $P_G(k) = P_{G_D}(k) - N(k)$ where

$$N(k) := \text{Number of ways coloring } G \text{ with } k \text{ colors such that } a, b \text{ are colored the same.} \quad (4)$$

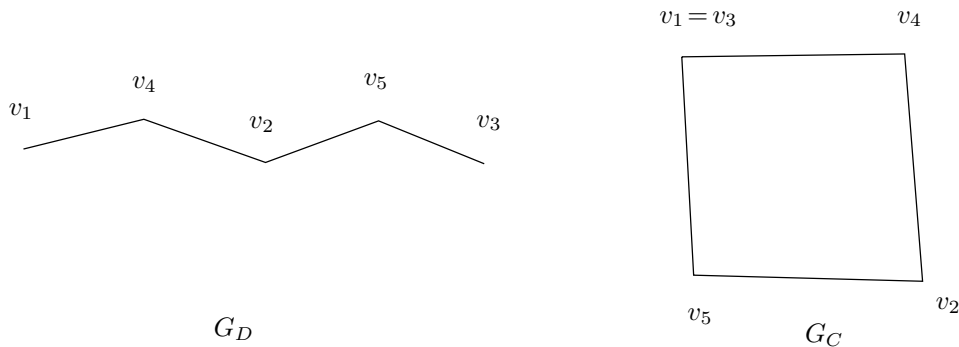
It is now clear that $N(k) = P_{G_C}(k)$. □

Example 9. We calculate $P_G(k)$ for the graph



through deletion-contraction. Take the edge $\{v_1, v_3\}$. It is easy to see that G_D is a chain, and G_C can be represented as a square cycle.

2. or “merging” a, b . Thus the edge e disappears. Also if both $\{c, a\}, \{c, b\} \in E$, then the two edges also “merge” into one edge in the graph G_C .



It is easy to see that

$$P_{G_D}(k) = k(k-1)^4. \tag{5}$$

On the other hand, we have

$$P_{G_C}(k) = k[(k-1)^2 + (k-1)(k-2)^2] = k(k-1)(k^2 - 3k + 3). \tag{6}$$

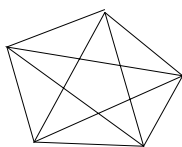
Thus we have

$$P_G(k) = P_{G_D}(k) - P_{G_C}(k) = k^5 - 5k^4 + 10k^3 - 10k^2 + 4k. \tag{7}$$

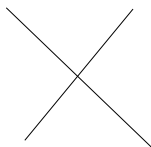
Exercise 2. Prove that the cycle with length n has $P_{C_n}(k) = (k-1)^n + (-1)^n(k-1)$. (Hint:³)

Exercise 3. Let W_n be the *wheel* of $n+1$ vertices, that is a cycle of n vertices with the $n+1$ -th vertex connected to them all. Find $P_{W_n}(k)$.

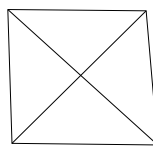
Exercise 4. Calculate $P_G(k)$ for the following graphs.



G_1



G_2



G_3

Answer: ⁴

Exercise 5. Prove the following “Two-Pieces Theorem”.

Theorem. (The Two-Pieces Theorem) Let the vertex set of G be partitioned into disjoint sets V_1, V_2 such that no edges in G joins a vertex in V_1 to a vertex in V_2 . Let G_1, G_2 be the subgraphs generated by V_1 and V_2 respectively. Then $P_G(k) = P_{G_1}(k)P_{G_2}(k)$.

Theorem 10. $P_G(k)$ is a polynomial of k with degree n and integer coefficients.

3. Induction.

4. $G_1: k^5 - 10k^4 + 35k^3 - 50k^2 + 24k$; $G_2: k^5 - 4k^4 + 6k^3 - 4k^2 + k$; $G_3: k^5 - 8k^4 + 24k^3 - 31k^2 + 14k$.

Proof. We prove through induction on the order n . When $n = 1$ the claim is clearly true. Now we assume that $P_G(k)$ is a polynomial of k with degree n and integer coefficients for every simple graph of order n .

Let $G = (V, E)$ be a simple graph of order $n + 1$. If $E = \emptyset$ then G is the null graph and $P_G(k) = k^{n+1}$. If $E \neq \emptyset$, let $e = \{a, b\} \in E$. By the deletion-contraction theorem we have

$$P_G(k) = P_{G_D}(k) - P_{G_C}(k). \quad (8)$$

Note that as G_C has order n , $P_{G_C}(k)$ is a polynomial of degree n . On the other hand, G_D is a simple graph of order $n + 1$ but with one less edge than G .

If G_D has no edges then $P_{G_D}(k) = k^{n+1}$ and the proof ends. If not, application of the deletion-contraction theorem to G_D yields

$$P_G(k) = P_{G_{DD}}(k) - P_{G_{DC}}(k) - P_{G_C}(k) \quad (9)$$

with $P_{G_{DC}}(k)$ also a degree n polynomial.

It is clear now that we can keep doing this until there is no edge left. At the end of the day we have

$$P_G(k) = P_{G_{D\dots D}}(k) - P_{G_{D\dots DC}}(k) - \dots - P_{G_{DC}}(k) - P_{G_C}(k) \quad (10)$$

with $P_{G_{D\dots D}}(k) = k^{n+1}$ and all the other polynomials of degree n .

Thus ends the proof. □

Exercise 6. Prove that the constant term of $P_G(k)$ is 0.

Exercise 7. Prove that the lead coefficient of $P_G(k)$ is always 1.

Exercise 8. Prove that unless $P_G(k) = k^n$, the sum of the coefficients is zero.

Exercise 9. Prove that the coefficient of k^{n-1} in $P_G(k)$ is the negative of the number of edges.

Exercise 10. Show that the following cannot be chromatic polynomials.

- a) $k^8 - 1$;
- b) $k^5 - k^3 + 2k$;
- c) $2k^3 - 3k^2$;
- d) $k^3 + k^2 + k$;
- e) $k^3 - k^2 + k$;
- f) $k^4 - 3k^3 + 3k^2$;
- g) $k^9 + k^8 - k^7 - k^6$.

Exercise 11. Prove that $P_G(k) \leq k(k-1)^{n-1}$ for any positive integer k , if G is connected with n vertices.

Exercise 12. Prove or disprove: If G' is a subgraph of G , then $P_{G'}(k) | P_G(k)$.

Exercise 13. Prove that if K_r is a subgraph of G , then $k(k-1)\dots(k-r+1) | P_G(k)$. (Hint:⁵)

5. Note that $P_G(i) = 0$ for $i = 0, 1, 2, \dots, r - 1$.