## On a Function Connected with $\varphi$ (n)

 $\mathbf{BY}$ 

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1. Let  $\varphi(n)$  denote Euler's function representing the number of numbers less than and prime to n. Let  $\varphi_1(n) = \varphi(n)$ ;  $\varphi_r(n) = \varphi(\varphi_{r-1}(n))$ ,  $r = 2, 3, \ldots$  For a given n,  $\varphi_r(n)$  decreases as r increases, and hence there is a least value of r, say  $r_1$ , such that

$$\varphi_r(n) = \varphi_{r+1}(n) = \ldots = 1.$$

S. S. Pillai proved that

$$\left[\frac{\log (n/2)}{\log 3}\right] + 1 \leqslant R(n) \leqslant \frac{\log n}{\log 2} + 1.$$

where  $r_1 = R(n)$ . In this paper I consider an analogous function S(n) given by

$$S(n) = \varphi_1(n) + \varphi_2(n) + ... + \varphi_{r_1}(n)$$

It is shown here that

$$R(n) \leqslant \frac{\log n}{\log 2} \text{ if } n \text{ is even}$$

and

$$\leq \frac{\log (n-1)}{\log 2} + 1$$
 if  $n$  is odd,

which is an improvement over Pillai's result;

Also

$$S(n) \leqslant n-1$$
 if  $n$  is even  $\leqslant 2n-3$  , odd,

while  $S(n) \ge 2[\log (n/2)/\log 3]+1-1$ .

I also consider the solutions of the equation S(n) = n, and show that each one of the following values of n provides a solution:

$$n = n_1, 3n_2, 3n_3, \ldots$$

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where 
$$n_1 = 3^k$$
,  $k = 0, 1, 2, ...$   
 $n_2 = 1 + 4n_1$ 

$$n_3 = 1 + 12n_2$$

$$n_4=1+12n_3,\ldots.$$

provided these are primes. The question whether these are the only primes remains open.

## 2. Theorem 1. If n is even,

$$R(n) \leqslant \frac{\log n}{\log 2}$$
;  $S(n) \leqslant n-1$ .

If 
$$n$$
 is odd,  $R(n) \leqslant \frac{\log (n-1)}{\log 2} + 1$ 

$$S(n) \leqslant 2n - 3$$
.

**Proof:** When n is even,  $\varphi_1(n) \leqslant (n/2)$ ;

$$\varphi_2(n) \leqslant \frac{1}{2} \varphi_1(n) \leqslant \frac{1}{4} n; \ldots$$

$$\varphi_r(n) \leqslant \frac{1}{2^r}n.$$

If  $r = r_1 = R(n)$  we get

$$1\leqslant \frac{1}{2^{r_1}}n,$$

$$r_1 \leqslant rac{\log \, n}{\log \, 2}$$

**(1)** 

If n is a power of 2,  $r_1$  is actually  $\left(\frac{\log n}{\log 2}\right)$ 

Again 
$$S(n) = \varphi_1(n) + \varphi_2(n) + ... + \varphi_{r1}^{(n)}$$

$$\leq \frac{1}{2}n + \frac{1}{4}n + \ldots + \frac{1}{2^{r_1}}n;$$

$$=n\left(1-rac{1}{2^{r_1}}
ight)$$

$$\leqslant n \left(1 - \frac{1}{n}\right)$$
, by (1)

Hence  $S(n) \leq n-1$ .

If n is a power of 2, it is easily seen that S(n) = n - 1.

Let n be odd. Then  $\varphi_2(n) \leqslant \frac{1}{2} \varphi_1(n)$ ;

$$\varphi_3(n)\leqslant \frac{1}{2}\,\varphi_2(n)\leqslant \frac{1}{2^2}\,\varphi_1(n)\,;\,\ldots.$$

$$\varphi_r(n) \leqslant \frac{1}{2^{r-1}} \varphi_1(n)$$

Hence

$$1 \leqslant \frac{1}{2^{r_1-1}} \varphi(n),$$

$$r_1 \leqslant \frac{\log \varphi(n)}{\log 2} + 1 \leqslant \frac{\log (n-1)}{\log 2} + 1 \qquad \qquad \dots$$
 (2)

If n is a prime of the form  $2^k + 1$ , we get the equality sign. Again

$$S(n) \leqslant \varphi(n) \left(1 + \frac{1}{2} + \dots + (1/2^{r_1 - 1})\right)$$

$$= 2\varphi(n) \left(1 - (1/2^{r_1})\right)$$

$$\leqslant 2(n - 1) \left(1 - (1/2^{r_1})\right)$$

$$\leqslant 2(n - 1) \left(1 - (1/2n - 1)\right), \text{ by } (2)$$

$$= 2n - 3.$$

Actually S(n) attains this value if n is a prime of the form  $2^k + 1$ .

3. Before considering the lower bounds of R(n) and S(n) we prove the following:

Theorem 2.\* For a given value of R(n), the maximum value of n is  $2 \cdot 3^{R(n)-1}$ .

The proof of this depends on the

Lemma.  $R(pn) > R(n) + (\log n/\log p)$ , if p is a prime > 3. If p = 3 it becomes an equality.

Proof by induction: Suppose the result is true for all primes q < p so that

$$R(qn) \geqslant R(n) + \frac{\log q}{\log 3} \qquad .. \quad (3)$$

\*This is equivalent to Pillai's theorem II; but the proof by induction given here is different from Pillai's.

for all n > 0 and all q < p. We will prove the same for p. Let  $p - 1 = 2^a q_1^{b_1} q_2^{b_2} \dots$ 

Then 
$$R(n) = 1 + R[\varphi(n)]$$
 .. (4)

assuming as we may that n > 2; for if n = 2, the result of the lemma is obvious.

Also 
$$R(pn) = 1 + R(\varphi(pn)).$$

Assume that p is prime to n as a first case. Then

$$egin{align} \mathrm{R}\left(pn
ight) &= 1 + \mathrm{R}\left[\left(p-1
ight) \phi\left(n
ight)
ight] \ &= 1 + \mathrm{R}\left[2^a q_1^{b_1} q_2^{b_2} \ldots \phi\left(n
ight)
ight] \ &= 1 + \mathrm{R}\left[q_1^{b_1} q_2^{b_2} \ldots \phi\left(n
ight)
ight] + a, \end{split}$$

since  $\varphi(n) \geqslant 2$  and R(2m) = R(m) + 1 if m is even;

Hence 
$$R(pn) \geqslant 1 + R[\varphi(n)] + a + \sum_{i=1}^{n} \left(\frac{\log qi}{\log 3}\right) bi$$
,

by using (3) repeatedly,

$$= R(n) + a + \left\{ \frac{\log (p-1)/2^{a}}{\log 3} \right\}$$

$$= R(n) + \frac{\log (3/2)^{a} (p-1)}{\log 3}$$

$$\geqslant R(n) + \left( \frac{\log p}{\log 3} \right)$$

if  $(3/2)^a(p-1) \geqslant p$  which is true since  $a \geqslant 1$ , and  $p \geqslant 3$ . Since the result is true for p=3, the lemma follows in the case when p is prime to n. Next let  $(p, n) \neq 1$ . Then in each of the pairs  $(\varphi_1(pn); \varphi_1(n)); (\varphi_2(pn); \varphi_2(n)); \ldots p$  occurs upto a stage, say in the first k pairs, and in the pair  $(\varphi_{k+1}(pn); \varphi_{k+1}(n))$  the first member contains p while the second does not.

Let us call  $\varphi_{k+1}(np) = pu$ ;  $\varphi_{k+1}(n) = u$ ; then (p, u) = 1. Hence by the above result,

$$R(pu) > R(u) + \left(\frac{\log p}{\log 3}\right) \qquad \qquad .. \quad (4)$$

But

$$R(pn) = k + 1 + R(pu)$$

$$R(n) = k + 1 + R(u)$$

Hence adding k+1 to both sides of (A) we get

$$R(pn) > R(n) + \left(\frac{\log p}{\log 3}\right), p > 3,$$

so that the lemma follows in this case also. To prove Theorem 2, let R(n) = t (fixed), and consider all possible values of n satisfying this. The greatest of such n's, say  $n_1$ , must not contain a prime factor p greater than 3, for then

$$R(n_1) = R\left(\frac{n_1}{p} \cdot p\right) \ge R\left(\frac{n_1}{p}\right) + \left(\frac{\log p}{\log 3}\right)$$

$$\ge R\left(\frac{n_1}{p}\right) + \left[\frac{\log p}{\log 3}\right] + 1$$

$$= R\left(\frac{n_1}{p} \cdot 3^k\right),$$

where  $K = \left\lceil \frac{\log p}{\log 3} \right\rceil + 1$ , since R(3u) = R(u) + 1 for all u.

But  $\frac{n_1}{p} \cdot 3^k > n_1$ . Hence  $n_1$  is not the largest solution.

Let now  $n_1 = 3u$ ; then t = R(3u) = R(u) + 1 so that R(u) = t - 1. Assuming that the theorem is true for all values of R(n) < t, it follows that the maximum value of u is  $2.3^{t-2}$ . Hence  $3u = 2.3^{t-1}$  and the theorem follows by induction. From the theorem it follows, as proved by Pillai, that

$$R(n) \geqslant \left[\frac{\log n - \log 2}{\log 3}\right] + 1. \qquad .. \quad (5)$$

Theorem 3.  $S(n) \ge [\log (n/2)/\log 3] + 1$ .

For 
$$S(n) = \varphi_1(n) + \varphi_2(n) + ... + \varphi_t(n)$$
;  $t = R(n)$ 

Now 
$$\varphi_t(n) = 1$$
;  $\varphi_{t-1}(n) = 2$ ;  $\varphi_{t-2}(n) \geqslant 2^2$ ;  $\varphi_{t-3}(n) \geqslant 2^3$ ; ...

Hence 
$$S(n) \ge 2^{t-1} + 2^{t-2} + \dots + 2 + 1$$
  
=  $2^t - 1$ .

The theorem now follows by using (5).

(4) S(n) as a function of R(n). We will show that  $2^{R(n)} - 1 \leq S(n) \leq 3^{R(n)}$ .

For a given R(n) = t, proceeding as above we get  $S(n) \ge 2^t - 1 = 2^{R(n)} - 1$ . When n is a power of 2, this becomes an equality. Again, by Theorem 2,

$$\varphi_t(n) = 1; \ \varphi_{t-1}(n) = 2; \ \varphi_{t-2}(n) \leq 2 \cdot 3; \ \varphi_{t-3}(n) \leq 2 \cdot 3^2 \dots$$

$$\varphi_1(n) \leq 2 \cdot 3^{t-1}$$

Hence for a given t,

$$S(n) \leq 2 \cdot 3^{t-1} + 2 \cdot 3^{t-2} + \dots + 2 \cdot 3 + 2 + 1$$
  
=  $3^{t}$ .

S(n) actually reaches this value when  $n = 2 \cdot 3^{t-1}$ .

4. Let us finally consider the solutions of the equation S(n) = n. By Theorem 1, n can only be odd. It is easily verified that two sets of solutions are  $n = 3^k$  (k = 0, 1, ...) and n = 3p where p is a prime = 4.  $3^k + 1$ , k = 1, 2, ...; we will prove

Teorem 4. If S(3p) = 3p then S((p-1)/4) = ((p-1)/4), where p is a prime > 3.

Proof. 
$$3p = S(3p) = 2(p-1) + S(2(p-1))$$
  
=  $2(p-1) + 2S(p-1) + 1$ ,

using 
$$\cdot S(2u) = 2S(u) + 1$$
 if u is even

$$= S(u)$$
 if  $u$  is odd.

Hence 
$$S(p-1) = \frac{p+1}{2}$$

Since 
$$p > 3$$
,  $\frac{p-1}{2}$  is even,

Hence 
$$S\left(\frac{p-1}{2}\right) = \left(\frac{p-1}{4}\right)$$

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Since S(n) is always odd,  $\frac{p-1}{4}$  must be odd, p=4k+1 (k odd);

$$\begin{split} &\frac{p-1}{2} = 2\,k,\; k \;\; \text{odd} \;\; \mathbb{S} \;\; \left( \begin{array}{c} \frac{p-1}{2} \end{array} \right) &= \mathbb{S} \;\; \left( 2^{\frac{p}{4}} \right) &= \mathbb{S}. \;\; \left( \begin{array}{c} \frac{p-1}{4} \end{array} \right) \\ &= \mathbb{S} \;\; \text{since} \; \frac{p-1}{4} \;\; \text{is odd}. \end{split}$$

Hence the theorem follows.

With the help of this result, we can derive an infinite sequence of classes of solutions of S(n) = n.

One is n = 3p, where  $\frac{p-1}{4} = 3^a$  or  $p = 4 \cdot 3^a + 1 = p_1$  say.

Another is n=3  $p_2$  where  $p_2$  is a prime given by  $\frac{p_2-1}{4}=3$   $p_1$  or  $p_2=12p_1+1$ .

This sequence of solutions can be continued indefinitely. It is conjectured that this sequence exhausts all solutions.

## REFERENCE

S. S. Pillai (1929) .. On a function connected with φ (n); Bulletin of the American Math. Society, pp. 837-841,